# Characterizing Alternating Groups by Their Codegree 

## Sets

Mallory Dolorfino, Luke Martin, Zachary Slonim, Yuxuan Sun, Yong Yang

Kalamazoo College, Gonzaga University, University of California Berkeley, Haverford College, Texas State University

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## Preliminary Definitions

## Definition 1: Complex Representation

A complex representation of a group $G$ is a homomorphism $\rho: G \rightarrow$ $\mathrm{GL}(n, \mathbb{C})$ from $G$ to $n \times n$ invertible matrices over $\mathbb{C}$. A representation is irreducible if it has no proper, nontrivial subrepresentation.

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## Definition 2: Character

The character of a complex representation $\rho$ is the map $\chi_{\rho}: g \mapsto$ $\operatorname{Tr}(\rho(g))$ where $\operatorname{Tr}(\rho(g))$ is the trace of the matrix $\rho(g)$. We let $\operatorname{Irr}(G)$ denote the set of irreducible characters of $G$.

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## Definition 3: Codegree

The codegree of a character $\chi$ of a group $G$ is defined as $\operatorname{cod}(\chi):=$ $\frac{|G: \operatorname{ker}(\chi)|}{\chi(1)}$, where $\operatorname{ker}(\chi):=\operatorname{ker}(\rho)$.

## Definition 4: Codegree Set

The codegree set of a group $G$ is defined as $\operatorname{cod}(G):=\{\operatorname{cod}(\chi) \mid \chi \in$ $\operatorname{Irr}(G)\}$.

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## Conjecture (Kourovka Notebook Problem 20.79)

If $H$ is a finite nonabelian simple group and $G$ is a finite group with $\operatorname{cod}(G)=\operatorname{cod}(H)$, then $G \cong H$.

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Recently, various authors have proved this conjecture for various simple groups such as $\operatorname{PSL}(2, q),{ }^{2} B_{2}\left(2^{2 f+1}\right), \operatorname{PSL}(3,4)$, and all the sporadic groups. In this project, we verify this conjecture for $H$ any alternating group $\mathrm{A}_{n}$ for $n \geq 5$.

## Proof Outline

- Suppose for contradiction that $G$ is a minimal counterexample to the conjecture. That is, $\operatorname{cod}(G)=\operatorname{cod}\left(\mathrm{A}_{n}\right)$, but $G \not \approx \mathrm{~A}_{n}$.


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- Suppose for contradiction that $G$ is a minimal counterexample to the conjecture. That is, $\operatorname{cod}(G)=\operatorname{cod}\left(\mathrm{A}_{n}\right)$, but $G \not \approx \mathrm{~A}_{n}$.
- Let $N$ be a maximal normal subgroup of $G$. Then $G / N$ is simple, and $\operatorname{cod}(G / N) \subseteq \operatorname{cod}(G)=\operatorname{cod}\left(\mathrm{A}_{n}\right)$.


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- We first show that $G / N \cong \mathrm{~A}_{n}$.
- Then, we show that $N=1$, so $G \cong G / N \cong \mathrm{~A}_{n}$.


## Proving $G / N \cong \mathrm{~A}_{n}$

We show that if $G / N \nsupseteq \mathrm{~A}_{n}$, we get a contradiction by proving $\operatorname{cod}(G / N) \nsubseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$ for all other nonabelian simple groups $G / N$.

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## Lemma 5: Dividing Orders

Let $H$ be a finite group and $K$ be a finite simple group with $\operatorname{cod}(K) \subseteq$ $\operatorname{cod}(H)$, then $|K|$ divides $|H|$.

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We show that if $G / N \neq \mathrm{A}_{n}$, we get a contradiction by proving $\operatorname{cod}(G / N) \nsubseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$ for all other nonabelian simple groups $G / N$.

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Proof. $\forall \chi \in \operatorname{Irr}(H), \chi(1)^{2}<|H|$. $H$ is simple so $\operatorname{ker}(\chi)=1$ and $\operatorname{cod}(\chi)=\frac{|H|}{\chi(1)}>\sqrt{|H|}$. Then, $\operatorname{cod}(K) \subseteq \operatorname{cod}(H)$ implies that $\forall \psi \in \operatorname{Irr}(K), \frac{|K|}{\psi(1)}>\operatorname{cod}(\psi)>\sqrt{|H|}$. Thus, $\psi(1)<\frac{|K|}{\sqrt{|H|}}$. Summing the squares of irreducible character degrees gives $|K|<|\operatorname{Irr}(K)| \frac{|K|^{2}}{|H|}$ and the inequality follows.

## $G / N$ is Sporadic

## Corollary 7: Order Restrictions

If $\operatorname{cod}(G / N) \subseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$, then $\frac{\left|\mathrm{A}_{n}\right|}{|G / N|} \in\{1,2, \ldots,|\operatorname{Irr}(G / N)|\}$.

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- For each sporadic (or Tits) group, $G / N$, we obtain an upper bound on $n$ by using $\left|A_{n}\right|<|G / N| \cdot|\operatorname{Irr}(G / N)|$.
- Then we can computationally check which values of $n$ satisfy $|G / N|$ divides $\left|\mathrm{A}_{n}\right|$.


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- The only possibility is $\mathrm{A}_{n}=\mathrm{A}_{10}$ and $G / N \cong \mathrm{~J}_{2}$,
- In this case, we can check using the ATLAS that $\operatorname{cod}\left(J_{2}\right) \nsubseteq \operatorname{cod}\left(\mathrm{A}_{10}\right)$.


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Thus if $\operatorname{cod}(G / N) \subseteq \operatorname{cod}\left(A_{n}\right), G / N$ cannot be a sporadic group or the Tits group.

## $G / N$ is a Simple Group of Lie Type

## Example 8: $\operatorname{PSL}(m+1, q)$

Let $G / N=\operatorname{PSL}(m+1, q)$, where $m$ is a positive integer and $q=p^{k}$ is a prime power. Then, $|G / N|=\frac{q^{m(m+1) / 2}}{\operatorname{gcd}(m+1, q-1)} \prod_{i=1}^{m}\left(q^{i+1}-1\right)$ and $|\operatorname{Irr}(G / N)| \leq 2.5 q^{m}$.

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Proof. We know $q^{m(m+1) / 2}$ divides $|G / N|$ and that $\left|\mathrm{A}_{n}\right|_{q} \leq \frac{n}{k(p-1)}$. Since $|G / N|$ divides $\left|A_{n}\right|$, we have $\frac{m(m+1)}{2} \leq \frac{n}{k(p-1)}$. Thus, $n \geq \frac{m(m+1) k(p-1)}{2}$.
Using the Lemma which limits $\left|A_{n}\right|$ and the above inequality, we have $\left|\mathrm{A}_{\frac{m(m+1) k(p-1)}{2}}\right|<|G / N| \cdot 2.5 q^{m}$. As functions of $m, p$, or $k$ only, the left-hand side grows faster asymptotically than the right so we find maximum values for each of these variables for which the inequality is satisfied. Namely, $m \leq 6, p \leq 17$, and $k \leq 63$. Now, we can anaylze each of these possible combinations in turn in the same way we did for the sporadic groups.

## $G / N$ is a Simple Group of Lie Type, Part 2

## Possible Exceptions

We get the following list of possible exceptions which satisfy $|\operatorname{PSL}(m+1, q)|$ divides $\left|\mathrm{A}_{n}\right|$ and $\left|\mathrm{A}_{n}\right|<|\operatorname{PSL}(m+1, q)| \cdot 2.5 q^{m}$ :

| $m$ | $q$ | $n$ |
| :---: | :---: | :---: |
| 1 | 4 | 5 |
| 1 | 4 | 6 |
| 1 | 8 | 7 |
| 1 | 9 | 6 |
| 1 | 9 | 7 |
| 1 | 5 | 5 |
| 1 | 5 | 6 |
| 1 | 7 | 7 |
| 2 | 4 | 8 |
| 2 | 4 | 9 |
| 3 | 2 | 8 |
| 3 | 2 | 9 |

Remark. $\operatorname{PSL}(2,4) \cong \operatorname{PSL}(2,5) \cong A_{5}$, $\operatorname{PSL}(2,9) \cong \mathrm{A}_{6}$, and $\operatorname{PSL}(4,2) \cong \mathrm{A}_{8}$.

- In all other cases, when $\operatorname{PSL}(m+1, q) \not \approx \mathrm{A}_{n}$, we can check using the ATLAS that $\operatorname{cod}(\operatorname{PSL}(m+1, q)) \nsubseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$.
- We follow these same steps to show that for $G / N$ any simple group of Lie type, $\operatorname{cod}(G / N) \nsubseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$.


## $G / N$ is an Alternating Group

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Let $a_{x}$ denote the the minimal nontrivial codegree of $A_{x}$. We will show that $a_{m}<a_{n}$, proving that if $\operatorname{cod}\left(\mathrm{A}_{m}\right) \subseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$, then $m=n$.

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Let $\lambda=\left(\lambda_{1} \geq \lambda_{2} \geq \ldots \geq \lambda_{k}\right)$ be a partition of $n$ and $V_{\lambda}$ be the corresponding irreducible representation of $S_{n}$.

## $G / N$ is an Alternating Group

Interpreting $\lambda$ as a Young diagram with $k$ rows of length $\lambda_{1} \ldots \lambda_{k}$, we have that

$$
\operatorname{cod}\left(V_{\lambda}\right)=\frac{n!}{\operatorname{dim}\left(V_{\lambda}\right)}=\prod h_{\lambda}(i, j):=H_{\lambda},
$$

where the product goes over all the cells of the Young diagram, and $h_{\lambda}(i, j)$ is the hook length of the cell $(i, j)$, i.e. the number of cells $(a, b)$ such that $a=i$ and $b \geq j$ or $a \geq i$ and $b=j$.

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Let $U_{\lambda}$ be the restriction of $V_{\lambda}$ to $A_{n}$. If $\lambda$ is not self-conjugate, then $U_{\lambda}$ remains irreducible, and, for our purposes, we can prove that we only need to consider $\lambda$ which are not self-conjugate.

## $G / N$ is an Alternating Group

Thus $a_{n}=H_{\lambda} / 2$ for some $\lambda$. Then for $m>2$, we may choose a non-self-conjugate partition of $m, \mu$, such that $\mu$ is completely contained $\lambda$. Then

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a_{m} \leq \frac{H_{\mu}}{2}<\frac{H_{\lambda}}{2}=a_{n}
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However, $a_{m} \in \operatorname{cod}\left(\mathrm{~A}_{m}\right) \subseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$ while $a_{n}$ is the smallest nontrivial element of $\operatorname{cod}\left(\mathrm{A}_{n}\right)$ and hence we get a contradiction. Thus if $G / N$ is an alternating group and $\operatorname{cod}(G / N) \subseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$, then $G / N \cong \mathrm{~A}_{n}$.

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Thus, we have shown in the previous slides that $\operatorname{cod}(G / N) \nsubseteq \operatorname{cod}\left(\mathrm{A}_{n}\right)$ for any simple group $G / N$ except $G / N \cong A_{n}$. Thus $G / N \cong A_{n}$.

## Proving $N=1$

Let $G$ be a minimal counterexample where $\operatorname{cod}(G)=\operatorname{cod}\left(A_{n}\right)$ but $G \not \approx \mathrm{~A}_{n}$. Let $N$ be a maximal normal subgroup of $G$. Then, $G / N \cong \mathrm{~A}_{n}$. We prove the following relatively simple properties:

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- $N$ is also a minimal normal subgroup
- $N$ is the unique minimal normal subgroup of $G$ which implies it is the only proper, nontrivial normal subgroup of $G$
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- $\mathbf{C}_{G}(N)=N$ (Schur multiplier of $\left.\mathrm{A}_{n}\right)$


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- $N$ is an elementary abelian group
- $\mathbf{C}_{G}(N)=N$ (Schur multiplier of $\mathrm{A}_{n}$ )
- $\frac{\left|I_{G}(\lambda)\right|}{\theta(1)} \in \operatorname{cod}(G)$ and $|N|$ divides $|G / N|$


## Proving $N=1$, Part 2

$$
\begin{aligned}
& \text { Summary } \\
& |N|=p^{m} \text { divides }|G / N|=\left|\mathrm{A}_{n}\right| \text {. By the normalizer-centralizer theorem, } \\
& \mathrm{A}_{n} \cong G / N=\mathbf{N}_{G}(N) / \mathbf{C}_{G}(N) \leq \operatorname{Aut}(N) \text { and } m>1 \text {. In general, } \\
& \operatorname{Aut}(N)=\operatorname{GL}(m, p) \text {, so } \mathrm{A}_{n} \lesssim \operatorname{GL}(m, p) \text {. }
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## Proving $N=1$, Part 2

## Summary <br> $|N|=p^{m}$ divides $|G / N|=\left|\mathrm{A}_{n}\right|$. By the normalizer-centralizer theorem, $\mathrm{A}_{n} \cong G / N=\mathbf{N}_{G}(N) / \mathbf{C}_{G}(N) \leq \operatorname{Aut}(N)$ and $m>1$. In general, $\operatorname{Aut}(N)=\operatorname{GL}(m, p)$, so $\mathrm{A}_{n} \lesssim \mathrm{GL}(m, p)$.

- The $p$-part of $n$ ! gives $\left|\mathrm{A}_{n}\right|_{p} \leq \frac{n}{p-1}$ for any prime $p$. I.E. $m \leq \frac{n}{p-1}$


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- For $n>7$, the minimal faithful representation of $\mathrm{A}_{n}$ over a finite field has degree at least $n-2$. I.E. $m \geq n-2$
- Thus, we deduce from these inequalities that $p=2$ and $m=n-2$


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- $\frac{|T|}{\vartheta(1)} \in \operatorname{cod}(G)$ and $|N|$ divides $\frac{|T|}{\vartheta(1)}$
- $\left|\frac{|T / N|}{\vartheta(1)}\right|_{2}=0$ so the 2-parts of $|T / N|$ and $\vartheta(1)$ are equal


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- $|N|_{2}=\left|\mathrm{A}_{n}\right|_{2}=m-2$
- Let $\lambda \in \operatorname{Irr}(N), \vartheta \in \operatorname{Irr}\left(I_{G}(\lambda) \mid \lambda\right)$, and $T:=I_{G}(\lambda)$
- $1<|G: T|<|N|=2^{n-2}$ for $|G: T|$ is the number of all conjugates of $\lambda$
- $\frac{|T|}{\vartheta(1)} \in \operatorname{cod}(G)$ and $|N|$ divides $\frac{|T|}{\vartheta(1)}$
- $\left|\frac{|T / N|}{\vartheta(1)}\right|_{2}=0$ so the 2-parts of $|T / N|$ and $\vartheta(1)$ are equal
- $|T / N|$ is a sum of squares of the form $\vartheta(1)$ so the 2-parts of each must be 1


## Proving $N=1$, Part 3

- $|N|_{2}=\left|\mathrm{A}_{n}\right|_{2}=m-2$
- Let $\lambda \in \operatorname{Irr}(N), \vartheta \in \operatorname{Irr}\left(I_{G}(\lambda) \mid \lambda\right)$, and $T:=I_{G}(\lambda)$
- $1<|G: T|<|N|=2^{n-2}$ for $|G: T|$ is the number of all conjugates of $\lambda$
- $\frac{|T|}{\vartheta(1)} \in \operatorname{cod}(G)$ and $|N|$ divides $\frac{|T|}{\vartheta(1)}$
- $\left|\frac{|T / N|}{\vartheta(1)}\right|_{2}=0$ so the 2-parts of $|T / N|$ and $\vartheta(1)$ are equal
- $|T / N|$ is a sum of squares of the form $\vartheta(1)$ so the 2-parts of each must be 1
- Thus, since $|G / N|_{2} \geq|N|_{2}=n-2$, we have $|G: T|_{2}=|G / N: T / N|_{2} \geq n-2$ and so $|G: T| \geq 2^{n-2}=|N|$, which is a contradiction.


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